

The Tightly Super 3-Extra Connectivity and Diagnosability of Locally Twisted Cubes

Mujiangshan Wang¹, Yunxia Ren², Yuqing Lin¹, Shiyong Wang²

¹School of Electrical Engineering and Computer Science, The University of Newcastle, Newcastle, NSW, Australia

²School of Mathematics and Information Science, Henan Normal University, Xinxiang, Henan, China

Email: wangshiyong@htu.edu.cn, shiyong@sxu.edu.cn

How to cite this paper: Wang, M., Ren, Y.X., Lin, Y.Q. and Wang, S.Y. (2017) The Tightly Super 3-Extra Connectivity and Diagnosability of Locally Twisted Cubes. *American Journal of Computational Mathematics*, 7, 127-144.

<https://doi.org/10.4236/ajcm.2017.72011>

Received: March 20, 2017

Accepted: June 5, 2017

Published: June 8, 2017

Copyright © 2017 by authors and Scientific Research Publishing Inc.

This work is licensed under the Creative Commons Attribution International License (CC BY 4.0).

<http://creativecommons.org/licenses/by/4.0/>



Open Access

Abstract

Diagnosability of a multiprocessor system G is one important measure of the reliability of interconnection networks. In 2016, Zhang *et al.* proposed the g -extra diagnosability of G , which restrains that every component of $G - S$ has at least $(g + 1)$ vertices. The locally twisted cube LTQ_n is applied widely. In this paper, we show that LTQ_n is tightly $(4n - 9)$ super 3-extra connected for $n \geq 6$ and the 3-extra diagnosability of LTQ_n under the PMC model and MM^* model is $4n - 6$ for $n \geq 5$ and $n \geq 7$, respectively.

Keywords

Interconnection Network, Combinatorics, Diagnosability

1. Introduction

At present, semiconductor technology has been widely applied in various fields of large-scale computer systems. But processors or communication links failures of a multiprocessor system give our live a lot of troubles. How to find out the faulty processors accurately and timely becomes the primary problem when the system is in operation. The diagnosis of the system is the process of identifying the faulty processors from the fault-free ones.

There are two well-known diagnosis models, one is the PMC diagnosis model, introduced by Preparata *et al.* [1] and the other is the MM model, proposed by Maeng and Malek [2]. In the PMC model, any two neighbor processors can test each other. In the MM model, to diagnose a system, we can compare their responses after a node sends the same task to its two neighbors. Sengupta and Dahbura [3] suggested a further modification of the MM model, called the MM^* model, in which each node must test another two neighbors.

In 1996, the g -extra connectivity $\tilde{\kappa}^{(g)}(G)$ of an interconnection network G

was introduced by Fàbrega and Fiol [4]. The g -extra connectivity $\tilde{\kappa}^{(g)}(G)$ of an interconnection network G has been widely studied [4]-[13].

In 2012, Peng *et al.* [14] proposed a measure for faulty diagnosis of the system, namely, the g -good-neighbor diagnosability, which restrains every fault-free node containing at least g fault-free neighbors. In [14], they studied the g -good-neighbor diagnosability of the n -dimensional hypercube under the PMC model. In 2016, Wang and Han [15] studied the g -good-neighbor diagnosability of the n -dimensional hypercube under the MM* model. In 2016, Zhang *et al.* [16] proposed the g -extra diagnosability of the system, which restrains that every component of $G - S$ has at least $(g + 1)$ vertices and showed the g -extra diagnosability of hypercubes under the PMC model and MM* model. Ren *et al.* [17] studied the tightly super 2-extra connectivity and 2-extra diagnosability of locally twisted cubes LTQ_n . In 2016, Wang *et al.* [18] studied the 2-extra diagnosability of the bubble-sort star graph BS_n under the PMC model and MM* model. In 2017, Wang and Yang [19] studied the 2-good-neighbor (2-extra) diagnosability of alternating group graph networks under the PMC model and MM* model.

In this paper, we show that LTQ_n is tightly $(4n - 9)$ super 3-extra connected for $n \geq 6$ and the 3-extra diagnosability of LTQ_n under the PMC model and MM* model is $4n - 6$ for $n \geq 5$ and $n \geq 7$, respectively.

2. Preliminaries

2.1. Notations

A multiprocessor system is modeled as an undirected simple graph $G = (V, E)$, whose vertices (nodes) represent processors and edges (links) represent communication links. Suppose that V' is a nonempty vertex subset of V . The induced subgraph by V' in G , denoted by $G[V']$, is a graph, whose vertex set is V' and whose edge set consists of all the edges of G with both endpoints in V' . The degree $d_G(v)$ of a vertex v in G is the number of edges incident with v . We denote by $\delta(G)$ the minimum degree of vertices of G . For any vertex v , we define the neighborhood $N_G(v)$ of v in G to be the set of vertices adjacent to v . u is called a neighbor vertex or a neighbor of v for $u \in N_G(v)$. Let $S \subseteq V(G)$. We denote by $N_G(S)$ the set $\bigcup_{v \in S} N_G(v) \setminus S$. For neighborhoods and degrees, we will usually omit the subscript for the graph when no confusion arises. A graph G is said to be k -regular if $d_G(v) = k$ for any vertex $v \in V$. A bipartite graph is one whose each edge has one end in subsets of vertex X and one end in subsets of vertex Y ; such a partition (X, Y) is called a bipartition of the graph. A complete bipartite graph is a simple bipartite graph with bipartition (X, Y) in which each vertex of X is joined to each vertex of Y ; if $|X| = m$ and $|Y| = n$, such a graph is denoted by $K_{m,n}$. The connectivity $\kappa(G)$ of a connected graph G is the minimum number of vertices whose removal results in a disconnected graph or only one vertex left. Let F_1 and F_2 be two distinct subsets of V , and let the symmetric difference $F_1 \Delta F_2 = (F_1 \setminus F_2) \cup (F_2 \setminus F_1)$. For graph-theoretical terminology and notation not defined here we follow [20].

Let $G=(V,E)$ be a connected graph. A faulty set $F \subseteq V$ is called a g -good-neighbor faulty set if $|N(v) \cap (V \setminus F)| \geq g$ for every vertex v in $V \setminus F$. A g -good-neighbor cut of G is a g -good-neighbor faulty set F such that $G - F$ is disconnected. The minimum cardinality of g -good-neighbor cuts is said to be the g -good-neighbor connectivity of G , denoted by $\kappa^{(g)}(G)$. A faulty set $F \subseteq V$ is called a g -extra faulty set if every component of $G - F$ has at least $(g+1)$ vertices. A g -extra cut of G is a g -extra faulty set F such that $G - F$ is disconnected. The minimum cardinality of g -extra cuts is said to be the g -extra connectivity of G , denoted by $\tilde{\kappa}^{(g)}(G)$.

Proposition 1. ([21]) *Let G be a g -extra and g -good-neighbor connected graph. Then $\tilde{\kappa}^{(g)}(G) \leq \kappa^{(g)}(G)$.*

Proposition 2. ([21]) *Let G be a 1-good-neighbor connected graph. Then $\kappa^{(1)}(G) = \tilde{\kappa}^{(1)}(G)$.*

2.2. Definitions and Propositions

Definition 3. ([22] [23] [24] [25]) *A system G is said to be t -diagnosable if all faulty processors can be identified without replacement, provided that the number of faults presented does not exceed t . The diagnosability of G is the maximum value of t such that G is t -diagnosable.*

For the PMC model and MM* model, we follow [26]. Under the PMC model, to diagnose a system $G=(V(G), E(G))$, two adjacent nodes in G are capable to perform tests on each other. For two adjacent nodes u and v in $V(G)$, the test performed by u on v is represented by the ordered pair (u, v) . The outcome of a test (u, v) is 1 (resp. 0) if u evaluate v as faulty (resp. fault-free). We assume that the testing result is reliable (resp. unreliable) if the node u is fault-free (resp. faulty). A test assignment T for G is a collection of tests for every adjacent pair of vertices. The collection of all test results for a test assignment T is called a syndrome. For a given syndrome σ , a subset of vertices $F \subseteq V(G)$ is said to be consistent with σ if syndrome σ can be produced from the situation that, for any $(u, v) \in L$ such that $u \in V \setminus F$, $\sigma(u, v) = 1$ if and only if $v \in F$. Let $\sigma(F)$ denote the set of all syndromes which F is consistent with. Under the PMC model, two distinct sets F_1 and F_2 in $V(G)$ are said to be indistinguishable if $\sigma(F_1) \cap \sigma(F_2) \neq \emptyset$, otherwise, F_1 and F_2 are said to be distinguishable.

Similar to the PMC model, we can define a subset of vertices $F \subseteq V(G)$ is consistent with a given syndrome σ^* and two distinct sets F_1 and F_2 in $V(G)$ are indistinguishable (resp. distinguishable) under the MM* model.

In a system $G=(V,E)$, a faulty set $F \subseteq V$ is called a g -extra faulty set if every component of $G - F$ has more than g nodes. G is g -extra t -diagnosable if and only if for each pair of distinct faulty g -extra vertex subsets $F_1, F_2 \subseteq V(G)$ such that $|F_i| \leq t$, F_1 and F_2 are distinguishable. The g -extra diagnosability of G , denoted by $\tilde{t}_g(G)$, is the maximum value of t such that G is g -extra t -diagnosable.

Proposition 4. [18] *For any given system G , $\tilde{t}_g(G) \leq \tilde{t}_{g'}(G)$ if $g \leq g'$.*

For an integer $n \geq 1$, a binary string of length n is denoted by $u_1 u_2 \cdots u_n$, where $u_i \in \{0, 1\}$ for any integer $i \in \{1, 2, \dots, n\}$. The n -dimensional locally twisted cube, denoted by LTQ_n , is an n -regular graph of 2^n vertices and $n2^{n-1}$ edges, which can be recursively defined as follows [27].

Definition 5. ([27]) For $n \geq 2$, an n -dimensional locally twisted cube, denoted by LTQ_n , is defined recursively as follows.

1) LTQ_2 is a graph consisting of four nodes labeled with 00, 01, 10 and 11, respectively, connected by four edges $\{00, 01\}$, $\{01, 11\}$, $\{11, 10\}$ and $\{10, 00\}$.

2) For $n \geq 3$, LTQ_n is built from two disjoint copies of LTQ_{n-1} according to the following steps. Let $0LTQ_{n-1}$ denote the graph obtained from one copy of LTQ_{n-1} by prefixing the label of each node with 0. Let $1LTQ_{n-1}$ denote the graph obtained from the other copy of LTQ_{n-1} by prefixing the label of each node with 1. Connect each node $0u_2 u_3 \cdots u_n$ of $0LTQ_{n-1}$ to the node $1(u_2 + u_n)u_3 \cdots u_n$ of $1LTQ_{n-1}$ with an edge, where “+” represents the modulo 2 addition.

The edges whose end vertices in different $iLTQ_{n-1}$ s are called to be cross-edges. **Figures 1-3** show four examples of locally twisted cubes. The locally twisted cube can also be equivalently defined in the following non-recursive fashion.

Definition 6. ([27]) For $n \geq 2$, the n -dimensional locally twisted cube, denoted by LTQ_n , is a graph with $\{0, 1\}^n$ as the node set. Two nodes $u_1 u_2 \cdots u_n$ and $v_1 v_2 \cdots v_n$ of LTQ_n are adjacent if and only if either one of the following conditions are satisfied.

1) $u_i = \bar{v}_i$ and $u_{i+1} = (v_{i+1} + v_n)(\text{mod } 2)$ for some $1 \leq i \leq n-2$, $n \geq 3$ and $u_j = v_j$ for all the remaining bits;

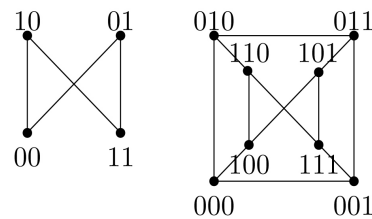


Figure 1. LTQ_2 and LTQ_3 .

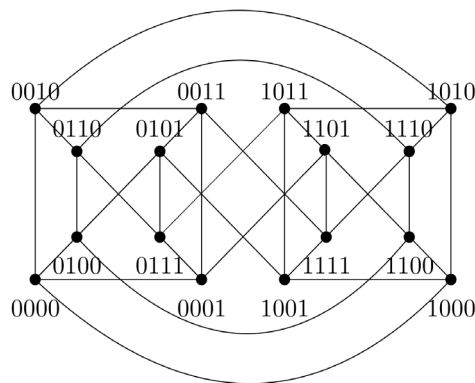
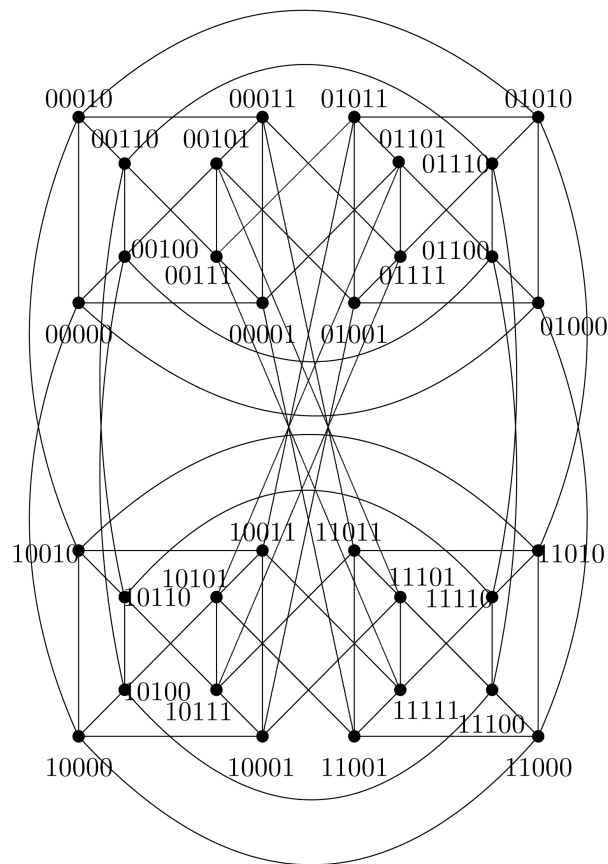


Figure 2. LTQ_4 .

Figure 3. LTQ_5 .

2) $u_i = \overline{v_i}$ for $i \in \{n-1, n\}$, $n \geq 2$ and $u_j = v_j$ for all the remaining bits.

Proposition 7. ([28]) Let LTQ_n be the locally twisted cube. If two vertices u, v are adjacent, there is no common neighbor vertex of these two vertices, i.e., $|N(u) \cap N(v)| = 0$. If two vertices u, v are not adjacent, there are at most two common neighbor vertices of these two vertices, i.e., $|N(u) \cap N(v)| \leq 2$.

3. The Connectivity of Locally Twisted Cubes

Lemma 1. ([27]) Let LTQ_n be the locally twisted cube. Then $\kappa(LTQ_n) = n$.

Lemma 2. ([29]) Let LTQ_n be the locally twisted cube, and let $S \subseteq V(LTQ_n)$ and $n \geq 3$. If $LTQ_n - S$ is disconnected and $n \leq |S| \leq 2n - 3$, then $LTQ_n - S$ has exactly two components, one is trivial and the other is nontrivial.

Lemma 3. ([17]) Let LTQ_n be the locally twisted cube. Then all cross-edges of LTQ_n is a perfect matching.

Lemma 4. ([30]) Let LTQ_n be the locally twisted cube. Then $\kappa^{(2)}(LTQ_n) = 4n - 8$.

Lemma 5. Let LTQ_n be the locally twisted cube. If $P = uvwx$ is a 3-path in LTQ_n and $ux \notin E(LTQ_n)$ for $n \geq 3$, $|N(V(P))| \geq 4n - 9$.

Proof. We decompose LTQ_n into $0LTQ_{n-1}$ and $1LTQ_{n-1}$. Then $0LTQ_{n-1}$ and $1LTQ_{n-1}$ are isomorphic to LTQ_{n-1} . Without loss of generality, we have the following cases.

Case 1. $u, x \in V(0LTQ_{n-1})$ and $v, w \in V(1LTQ_{n-1})$.

Since $u \in V(0LTQ_{n-1})$, $v \in V(1LTQ_{n-1})$ and u, v are adjacent, by Proposition 7, u, v have no the common neighbor vertex. Similarly, x, w have no the common neighbor vertex and v, w have no the common neighbor vertex. Since $u \in V(0LTQ_{n-1})$, $w \in V(1LTQ_{n-1})$, u, w are not adjacent, v is a common neighbor vertex of u, w , $x \in V(0LTQ_{n-1})$ and x is a neighbor vertex of w , by Lemma 3, $|N(u) \cap N(w) \setminus \{v\}| = 0$. Similarly, $|N(x) \cap N(v) \setminus \{w\}| = 0$. Since u and x are not adjacent, by proposition 7, $|N(u) \cap N(x)| \leq 2$. Therefore, $|N(V(P))| \geq 2(n-1) + 2(n-2) - 2 = 4n - 8$.

Case 2. $u \in V(0LTQ_{n-1})$ and $v, w, x \in V(1LTQ_{n-1})$.

Since u, v are adjacent, by Proposition 7, $|N(u) \cap N(v)| = 0$. Similarly, $|N(v) \cap N(w)| = 0$, $|N(x) \cap N(w)| = 0$. And since $u \in V(0LTQ_{n-1})$, $w \in V(1LTQ_{n-1})$, u, w are not adjacent and v is the common neighbor vertex of u and w , by Lemma 3, $|N(u) \cap N(w) \setminus \{v\}| \leq 1$. Since u, x are not adjacent, $u \in V(0LTQ_{n-1})$, $x \in V(1LTQ_{n-1})$, by Lemma 3, $|N(u) \cap N(x)| \leq 1$. Since w is the common neighbor vertex of v and x and v, x are not adjacent, by proposition 7, $|N(v) \cap N(x) \setminus \{w\}| \leq 1$. Therefore, $|N(P)| \geq 2(n-1) + 2(n-2) - 3 = 4n - 9$.

Case 3. $u, v \in V(0LTQ_{n-1})$ and $w, x \in V(1LTQ_{n-1})$.

Since u, v are adjacent, by Proposition 7, $|N(u) \cap N(v)| = 0$. Similarly, $|N(u) \cap N(w)| = 0$, $|N(w) \cap N(x)| = 0$. Since $u \in V(0LTQ_{n-1})$, $x \in V(1LTQ_{n-1})$ and u, x are not adjacent, by proposition 7, $|N(u) \cap N(x)| \leq 2$. If $|N(u) \cap N(w) \setminus \{v\}| = 1$, then, by Lemma 3, $|N(u) \cap N(x)| \leq 1$. If $|N(u) \cap N(w) \setminus \{v\}| = 0$, then, by Lemma 3, $|N(u) \cap N(x)| \leq 2$. Therefore, $|N(V(P))| \geq 2(n-1) + 2(n-2) - 2 = 4n - 8$.

Case 4. $u, v, w, x \in V(1LTQ_{n-1})$.

This case is clear.

In conclusion, $|N(V(P))| \geq 4n - 9$.

Lemma 6. Let LTQ_n be the locally twisted cube. If $LTQ_n[\{u, v, w, x\}]$ is isomorphic to $K_{1,3}$ for $n \geq 3$ and $d(u) = 3$, then

$$|N(V(LTQ_n[\{u, v, w, x\}]))| \geq 4n - 9.$$

Proof. Since $d(u) = 3$ and $LTQ_n[\{u, v, w, x\}]$ is isomorphic to $K_{1,3}$, we have $d(v) = 1$, $d(w) = 1$ and $d(x) = 1$. Since v, w are not adjacent and u is a common neighbor vertex of v, w , by Proposition 7, $|N(v) \cap N(w) \setminus \{u\}| \leq 1$. Similarly, $|N(v) \cap N(x) \setminus \{u\}| \leq 1$, $|N(w) \cap N(x) \setminus \{u\}| \leq 1$. Therefore, $|N(V(LTQ_n[\{u, v, w, x\}]))| \geq 3(n-1) + (n-3) - 3 = 4n - 9$.

If $LTQ_n[\{u, v, w, x\}]$ is a 4-cycle, then $|N(V(LTQ_n[\{u, v, w, x\}]))| = 4n - 8$. Combining this with Lemmas 5 and 6, we have the following corollary.

Corollary 1. Let LTQ_n be the locally twisted cube and let H be a connected subgraph of LTQ_n . If $|V(H)| \geq 4$, then $|N(V(H))| \geq 4n - 9$.

Lemma 7. Let $A = \{0 \cdots 0001, 0 \cdots 0111, 0 \cdots 0101, 0 \cdots 0100\}$ and let LTQ_n be the locally twisted cube with $n \geq 4$. If $F_1 = N_{LTQ_n}(A)$, $F_2 = F_1 \cup A$, where $n \geq 4$, then $|F_1| = 4n - 9$, $|F_2| = 4n - 5$, F_1 is a 3-extra cut of LTQ_n , $LTQ_n - F_1$ has two components $LTQ_n - F_2$ and $LTQ_n[A]$,

$|V(LTQ_n - F_2)| \geq 4$, and $|A| \geq 4$.

Proof. According to the definition, $LTQ_n[A]$ is a 3-path and $|A| = 4$. By Lemma 5, $|F_1| \geq 4n - 9$. From Figure 2 and the definition of LTQ_n , we have that $|F_1| = 2(n - 3) + 2(n - 2) - 3 + 4 = 4n - 9$. Therefore,

$|F_2| = |F_1| + |A| = (4n - 9) + 4 = 4n - 5$. Let $F_2^i = V(iLTQ_{n-1}) \cap F_2$, $i \in \{0, 1\}$.

To prove $LTQ_n - F_2$ has two components and $|V(LTQ_n - F_2)| \geq 4$, we have the following discussion.

Claim 1. $LTQ_n - F_2$ is connected for $n \geq 4$.

The proof is by induction on n . For $n = 4$, $A = \{0001, 0111, 0101, 0100\}$, $F_1 = \{0000, 0011, 0110, 1001, 1011, 1101, 1100\}$. It is easy to see that $LTQ_4 - F_2$ is connected (See Figure 2). When $n = 5$, $A = \{00001, 00111, 00101, 00100\}$, $F_2^1 = \{11001, 11110, 11111, 10100\}$ (See Figure 3). It is clear that $1LTQ_{n-1} - F_2^1$ is connected (See Figure 3). We decompose LTQ_n into $0LTQ_{n-1}$ and $1LTQ_{n-1}$. Assume that $n \geq 6$, the result holds for LTQ_{n-1} . Then $0LTQ_{n-1} - F_2^0$ is connected. Note that $A \subseteq V(0LTQ_{n-1})$ and $|N(A) \cap V(1LTQ_{n-1})| = 4$. By Lemma 1, $1LTQ_{n-1} - F_2^1$ is connected. By inductive hypothesis, $0LTQ_{n-1} - F_2^0$ is connected. Since $2^{n-1} > 4n - 5$, by Lemma 3, $LTQ_n - F_2$ is connected. The proof of Claim 1 is complete.

By Claim 1, $LTQ_n - F_1$ has two components $LTQ_n - F_2$ and $LTQ_n[A]$ for $n \geq 4$. Then $|V(LTQ_n - F_2)| = 2^n - (4n - 5) \geq 4$ for $n \geq 4$. And since $|A| = 4$, F_1 is a 3-extra cut of LTQ_n .

Lemma 8. ([17]) Let LTQ_n ($n \geq 4$) be the locally twisted cube. If $|F| \leq 3n - 6$, then $LTQ_n - F$ satisfies one of the following conditions:

- 1) $LTQ_n - F$ has three components, two of which are isolated vertices;
- 2) $LTQ_n - F$ has two components, one of which is an isolated vertex;
- 3) $LTQ_n - F$ has two components, one of which is a K_2 ;
- 4) $LTQ_n - F$ is connected.

Theorem 8. ([31]) Let LTQ_n be the locally twisted cube. Then $\tilde{\kappa}^{(3)}(LTQ_n) = 4n - 9$ for $n \geq 4$.

Lemma 9. Let LTQ_n be the locally twisted cube. If $|F| = 10$ for $n = 5$, then $LTQ_5 - F$ satisfies one of the following conditions:

- 1) $LTQ_5 - F$ has four components, three of which are isolated vertices;
- 2) $LTQ_5 - F$ has three components, one of which is isolated vertices and one of which is a K_2 ;
- 3) $LTQ_5 - F$ has three components, two of which are isolated vertices;
- 4) $LTQ_5 - F$ has two components, one of which is a path of length two;
- 5) $LTQ_5 - F$ has two components, one of which is an isolated vertex;
- 6) $LTQ_5 - F$ has two components, one of which is a K_2 ;
- 7) $LTQ_5 - F$ is connected.

Proof. We decompose LTQ_5 into $0LTQ_4$ and $1LTQ_4$. Then $0LTQ_4$ and $1LTQ_4$ are isomorphic to LTQ_4 . Suppose that $F_i = F \cap V(iLTQ_4)$, $i \in \{0, 1\}$. Without loss of generality, let $|F_0| \geq |F_1|$. And since $|F| = 10$, $5 \leq |F_0| \leq 10$, $0 \leq |F_1| \leq 5$. Let C_i be the maximum component of $iLTQ_4 - F_i$, $i \in \{0, 1\}$. We consider the following cases.

Case 1. $|F_0| = 5$.

Since $|F_0| = 5$ and $|F| = 10$, $|F_1| = 10 - 5 = 5$. By Lemmas 1 and 2, both $0LTQ_4 - F_0$ and $1LTQ_4 - F_1$ are connected or has two components, one of which is an isolated vertex. Since $2^{5-1} - 6 - 2 \geq 1$, by Lemma 3,

$LTQ_n[V(C_0) \cup V(C_1)]$ is connected. Thus, $LTQ_5 - F$ satisfies one of conditions:

- 1) $LTQ_5 - F$ has three components, two of which are isolated vertices;
- 2) $LTQ_5 - F$ has two components, one of which is an isolated vertex;
- 3) $LTQ_5 - F$ has two components, one of which is a K_2 ;
- 4) $LTQ_5 - F$ is connected.

Case 2. $|F_0| = 6$.

Since $|F_0| = 6$ and $|F| = 10$, $|F_1| = 10 - 6 = 4$. By Lemmas 1 and 2, $1LTQ_4 - F_1$ is connected or has two components, one of which is an isolated vertex. Since $|F_0| = 6$, by Lemma 8, $0LTQ_4 - F_0$ satisfies one of the following conditions:

- 1) $0LTQ_4 - F_0$ has three components, two of which are isolated vertices;
- 2) $0LTQ_4 - F_0$ has two components, one of which is an isolated vertex;
- 3) $0LTQ_4 - F_0$ has two components, one of which is a K_2 ;
- 4) $0LTQ_4 - F_0$ is connected.

Then $LTQ_5 - F$ satisfies one of the conditions (1)-(7).

Case 3. $|F_0| \geq 7$.

Since $|F_0| \geq 7$ and $|F| = 10$, $|F_1| \leq 10 - 7 = 3$. By Lemma 1, $1LTQ_4 - F_1$ is connected.

Suppose that $0LTQ_4 - F_0$ is connected. Since $2^{5-1} - 10 \geq 1$, by Lemma 3, $LTQ_n - F$ is connected.

Suppose that $0LTQ_4 - F_0$ is not connected. Let the components in $0LTQ_4 - F_0$ be $G_1, G_2, \dots, G_k$ for $k \geq 2$ and $|V(G_1)| \leq |V(G_2)| \leq \dots \leq |V(G_k)|$. If $|V(G_r)| \geq 4$ ($1 \leq r \leq k-1$), by Lemma 3, $|N(V(G_r)) \cap V(1LTQ_4)| \geq 4$. Combining this with $|F_1| \leq 3$, we have that $LTQ_5[V(G_r) \cup V(1LTQ_4 - F_1)]$ is connected. Therefore, G_r is not a component of $LTQ_5 - F$ for $|V(G_r)| \geq 4$. Therefore, $LTQ_5 - F$ is connected. The following we discuss G_r is a component of $LTQ_5 - F$ with $|V(G_r)| \leq 3$ ($1 \leq r \leq k-1$).

If $k = 5$, by Lemma 3,

$|N(V(G_1)) \cup N(V(G_2)) \cup \dots \cup N(V(G_{k-1})) \cap V(1LTQ_4)| \geq 4$. Combining this with $|F_1| \leq 3$, there is one G_r ($1 \leq r \leq k-1$) such that

$LTQ_5[V(G_r) \cup V(1LTQ_4 - F_1)]$ is connected. Thus, $k \leq 4$. Since $|F_1| = 3$, $k \leq 4$, and $|V(G_r)| \leq 3$ ($1 \leq r \leq k-1$), $LTQ_5 - F$ satisfies one of the conditions (1)-(7).

Lemma 10. Let LTQ_n be the locally twisted cube. If $3n - 5 \leq |F| \leq 4n - 10$ for $n \geq 5$, then $LTQ_n - F$ satisfies one of the following conditions:

- 1) $LTQ_n - F$ has four components, three of which are isolated vertices;
- 2) $LTQ_n - F$ has three components, one of which is isolated vertices and one of which is a K_2 ;
- 3) $LTQ_n - F$ has three components, two of which are isolated vertices;

- 4) $LTQ_n - F$ has two components, one of which is a path of length two;
- 5) $LTQ_n - F$ has two components, one of which is an isolated vertex;
- 6) $LTQ_n - F$ has two components, one of which is a K_2 ;
- 7) $LTQ_n - F$ is connected.

Proof. By Lemma 9, the result holds for $n = 5$. We proceed by induction on n . Assume $n \geq 6$ and the result holds for LTQ_{n-1} , i.e., if $3n - 5 \leq |F| \leq 4(n-1) - 10 = 4n - 14$, then $LTQ_{n-1} - F$ satisfies one of the conditions (1)-(7) in Lemma 10. The following we prove $LTQ_n - F$ satisfies one of the conditions (1)-(7).

We decompose LTQ_n into $0LTQ_{n-1}$ and $1LTQ_{n-1}$. Then $0LTQ_{n-1}$ and $1LTQ_{n-1}$ are isomorphic to LTQ_{n-1} . Suppose that $F_i = F \cap V(iLTQ_{n-1})$, $i \in \{0, 1\}$. Without loss of generality, let $|F_0| \geq |F_1|$. And since

$$3n - 5 \leq |F| \leq 4n - 10, \quad n \leq \left\lceil \frac{3n-5}{2} \right\rceil \leq |F_0| \leq 4n - 10, \quad 0 \leq |F_1| \leq \left\lfloor \frac{4n-10}{2} \right\rfloor \leq 2n - 5.$$

Let C_i be the maximum component of $iLTQ_{n-1} - F_i$, $i \in \{0, 1\}$. We consider the following cases.

Case 1. $n \leq |F_0| \leq 3(n-1) - 6 = 3n - 9$.

Since $|F_0| \geq |F_1|$ and $|F| \leq 4n - 10$,

$$(4n - 10) - (3n - 9) = n - 1 \leq |F_1| \leq \left\lfloor \frac{4n-10}{2} \right\rfloor = 2n - 5. \text{ By Lemmas 1 and 2,}$$

$1LTQ_{n-1} - F_1$ is connected or has two components, one of which is an isolated vertex. Since $n \leq |F_0| \leq 3(n-1) - 6 = 3n - 9$, by lemma 8, $0LTQ_{n-1} - F_0$ satisfies one of the following conditions: 1) $0LTQ_{n-1} - F_0$ has three components, two of which are isolated vertices; 2) $0LTQ_{n-1} - F_0$ has two components, one of which is an isolated vertex; 3) $0LTQ_{n-1} - F_0$ has two components, one of which is a K_2 ; 4) $0LTQ_{n-1} - F_0$ is connected. Since $2^{n-1} - (4n - 10) - 3 \geq 1$, by Lemma 3, $LTQ_n[V(C_0) \cup V(C_1)]$ is connected. Thus, $LTQ_n - F$ satisfies one of conditions (1)-(7) in Lemma 10.

Case 2. $3n - 8 \leq |F_0| \leq 4n - 14$.

Since $|F_0| \geq |F_1|$ and $|F| \leq 4n - 10$, $|F_1| \leq (4n - 10) - (3n - 8) = n - 2$. By Lemma 1, $1LTQ_{n-1} - F_1$ is connected. Since $3n - 8 \leq |F_0| \leq 4n - 14$, according to inductive hypothesis, $0LTQ_{n-1} - F_0$ satisfies one of the following conditions:

- 1) $0LTQ_{n-1} - F_0$ has four components, three of which are isolated vertices;
- 2) $0LTQ_{n-1} - F_0$ has three components, one of which is isolated vertices and one of which is a K_2 ;
- 3) $0LTQ_{n-1} - F_0$ has three components, two of which are isolated vertices;
- 4) $0LTQ_{n-1} - F_0$ has two components, one of which is a path of length two;
- 5) $0LTQ_{n-1} - F_0$ has two components, one of which is an isolated vertex;
- 6) $0LTQ_{n-1} - F_0$ has two components, one of which is a K_2 ;
- 7) $0LTQ_{n-1} - F_0$ is connected.

Thus, $LTQ_n - F$ satisfies one of the conditions (1)-(7) in Lemma 10.

Case 3. $4n - 13 \leq |F_0| \leq 4n - 10$.

Since $4n - 13 \leq |F_0| \leq 4n - 10$ and $|F| \leq 4n - 10$,

$$|F_1| \leq (4n - 10) - (4n - 13) = 3. \text{ By Lemma 1, } 1LTQ_{n-1} - F_1 \text{ is connected.}$$

Suppose that $0LTQ_{n-1} - F_0$ is connected. Since $2^{n-1} - (4n-10) \geq 1$, by Lemma 3, $LTQ_n - F$ is connected.

Suppose that $0LTQ_{n-1} - F_0$ is not connected. Let the components in $0LTQ_{n-1} - F_0$ be $G_1, G_2, \dots, G_k$ for $k \geq 2$ and $|V(G_1)| \leq |V(G_2)| \leq \dots \leq |V(G_k)|$. If $|V(G_r)| \geq 4$ ($1 \leq r \leq k-1$), by Lemma 3, $|N(V(G_r)) \cap V(1LTQ_{n-1})| \geq 4$. Combining this with $|F_1| \leq (4n-10) - (4n-13) = 3$, we have that $LTQ_n[V(G_r) \cup V(1LTQ_{n-1} - F_1)]$ is connected. Therefore, G_r is not a component of $LTQ_n - F$ for $|V(G_r)| \geq 4$. Therefore, $LTQ_n - F$ is connected. The following we discuss G_r is a component of $LTQ_n - F$ with $|V(G_r)| \leq 3$ ($1 \leq r \leq k-1$).

If $k = 5$, by Lemma 3, $|N(V(G_1)) \cup N(V(G_2)) \cup \dots \cup N(V(G_{k-1})) \cap V(1LTQ_{n-1})| \geq 4$. Combining this with $|F_1| \leq 3$, there is one G_r ($1 \leq r \leq k-1$) such that $LTQ_n[V(G_r) \cup V(1LTQ_{n-1} - F_1)]$ is connected. Thus, $k \leq 4$. Since $|F_1| \leq 3$, $|V(G_r)| \leq 3$ ($1 \leq r \leq k-1$) and $k \leq 4$, $LTQ_n - F$ satisfies one of the conditions (1)-(7).

A connected graph G is super g -extra connected if every minimum g -extra cut F of G isolates one connected subgraph of order $g+1$. In addition, if $G - F$ has two components, one of which is the connected subgraph of order $g+1$, then G is tightly $|F|$ super g -extra connected.

Theorem 9. Let LTQ_n be the locally twisted cube for $n \geq 6$. Then LTQ_n is tightly $(4n-9)$ super 3-extra connected.

Proof. By Theorem 8, we know for any minimum 3-extra cut $F \subset V(LTQ_n)$, $|F| = 4n-9$. We decompose LTQ_n into $0LTQ_{n-1}$ and $1LTQ_{n-1}$. Then $0LTQ_{n-1}$ and $1LTQ_{n-1}$ are isomorphic to LTQ_{n-1} . Suppose that $F_i = F \cap V(iLTQ_{n-1})$, $i \in \{0, 1\}$. Without loss of generality, let $|F_0| \geq |F_1|$. And since $|F| = 4n-9$, $2n-4 \leq \left\lfloor \frac{4n-9}{2} \right\rfloor \leq |F_0| \leq 4n-9$, $0 \leq |F_1| \leq \left\lfloor \frac{4n-9}{2} \right\rfloor \leq 2n-5$. Let C_i be the maximum component of $iLTQ_{n-1} - F_i$, $i \in \{0, 1\}$. We consider the following cases.

Case 1. $2n-4 \leq |F_0| \leq 3(n-1)-6 = 3n-9$.

Since $|F_0| \geq |F_1|$ and $|F| = 4n-9$, $|F_1| \leq 2n-5$ holds.

By Lemmas 1 and 2, $1LTQ_{n-1} - F_1$ is connected or has two components, one of which is an isolated vertex. Since $2n-4 \leq |F_0| \leq 3(n-1)-6 = 3n-9$, by lemma 8, $0LTQ_{n-1} - F_0$ satisfies one of the following conditions: 1) $0LTQ_{n-1} - F_0$ has three components, two of which are isolated vertices; 2) $0LTQ_{n-1} - F_0$ has two components, one of which is an isolated vertex; 3) $0LTQ_{n-1} - F_0$ has two components, one of which is a K_2 ; 4) $0LTQ_{n-1} - F_0$ is connected. Since $2^{n-1} - (4n-9) - 3 \geq 1$, by Lemma 3, $LTQ_n[V(C_0) \cup V(C_1)]$ is connected. Then $LTQ_n - F$ satisfies one of the following conditions:

- 1) $LTQ_n - F$ has four components, three of which are isolated vertices;
- 2) $LTQ_n - F$ has three components, one of which is isolated vertices and one of which is a K_2 ;
- 3) $LTQ_n - F$ has three components, two of which are isolated vertices;

- 4) $LTQ_n - F$ has two components, one of which is a path of length two;
- 5) $LTQ_n - F$ has two components, one of which is an isolated vertex;
- 6) $LTQ_n - F$ has two components, one of which is a K_2 ;
- 7) $LTQ_n - F$ is connected.

Thus, in this case, F is not a minimum 3-extra cut of LTQ_n , a contradiction.

Case 2. $|F_0| = 3n - 8$.

Since $|F_0| = 3n - 8$ and $|F| = 4n - 9$, we have $|F_1| = (4n - 9) - (3n - 8) = n - 1$. By Lemmas 1 and 2, $1LTQ_{n-1} - F_1$ is connected or has two components, one of which is an isolated vertex. Since $|F_0| = 3n - 8$, by Lemma 10, $0LTQ_{n-1} - F_0$ satisfies one of the following conditions:

- 1) $0LTQ_{n-1} - F_0$ has four components, three of which are isolated vertices;
- 2) $0LTQ_{n-1} - F_0$ has three components, one of which is isolated vertices and the other of which is a K_2 ;
- 3) $0LTQ_{n-1} - F_0$ has three components, two of which are isolated vertices;
- 4) $0LTQ_{n-1} - F_0$ has two components, one of which is a path of length two;
- 5) $0LTQ_{n-1} - F_0$ has two components, one of which is an isolated vertex;
- 6) $0LTQ_{n-1} - F_0$ has two components, one of which is a K_2 ;
- 7) $0LTQ_{n-1} - F_0$ is connected.

If $0LTQ_{n-1} - F_0$ satisfies the condition (4), i.e., $0LTQ_{n-1} - F_0$ has two components, one of which is a path of length two, denoted by $P = uvw$, $1LTQ_{n-1} - F_1$ has two components, one of which is an isolated vertex x , and $|N(x) \cap V(P)| = 1$, $(N(V(P)) \cap V(1LTQ_{n-1})) \setminus \{x\} \subseteq F_1$, then, by Lemma 3, $LTQ_n - F$ has one component which is a 3-path or a $K_{1,3}$. Since $2^{n-1} - (4n - 9) - 3 \geq 1$ for $n \geq 6$, $LTQ_n[C_0 \cup C_1]$ is connected. Thus, $LTQ_n - F$ exactly has two components. Then the other component C satisfies $|C| = 2^n - (4n - 9) - 4 > 4$ for $n \geq 6$. Otherwise, F is not a minimum 3-extra cut of LTQ_n , a contradiction.

Case 3. $3n - 7 \leq |F_0| \leq 4n - 14$.

Since $|F_0| \geq |F_1|$ and $|F| \leq 4n - 9$, $|F_1| \leq (4n - 9) - (3n - 7) = n - 2$. By Lemma 1, $1LTQ_{n-1} - F_1$ is connected. Since $3n - 7 \leq |F_0| \leq 4n - 14$, by Lemma 10, $0LTQ_{n-1} - F_0$ satisfies one of the following conditions:

- 1) $0LTQ_{n-1} - F_0$ has four components, three of which are isolated vertices;
- 2) $0LTQ_{n-1} - F_0$ has three components, one of which is isolated vertices and the other of which is a K_2 ;
- 3) $0LTQ_{n-1} - F_0$ has three components, two of which are isolated vertices;
- 4) $0LTQ_{n-1} - F_0$ has two components, one of which is a path of length two;
- 5) $0LTQ_{n-1} - F_0$ has two components, one of which is an isolated vertex;
- 6) $0LTQ_{n-1} - F_0$ has two components, one of which is a K_2 ;
- 7) $0LTQ_{n-1} - F_0$ is connected.

Thus, $LTQ_n - F$ satisfies one of the following conditions:

- 1) $LTQ_n - F$ has four components, three of which are isolated vertices;
- 2) $LTQ_n - F$ has three components, one of which is isolated vertices and one of which is a K_2 ;
- 3) $LTQ_n - F$ has three components, two of which are isolated vertices;
- 4) $LTQ_n - F$ has two components, one of which is a path of length two;
- 5) $LTQ_n - F$ has two components, one of which is an isolated vertex;

- 6) $LTQ_n - F$ has two components, one of which is a K_2 ;
- 7) $LTQ_n - F$ is connected.

In this case, F is not a minimum 3-extra cut of LTQ_n , a contradiction.

Case 4. $|F_0| = 4n - 13$.

Since $|F_0| = 4n - 13$ and $|F| = 4n - 9$ for $n \geq 6$,

$|F_1| = (4n - 9) - (4n - 13) = 4$. By Lemma 1, $1LTQ_{n-1} - F_1$ is connected.

If there exists a 3-path P in $0LTQ_{n-1} - F_0$, then $N(V(P)) \cap V(0LTQ_{n-1}) \subseteq F_0$. By Corollary 1, $|N(V(P))| \geq 4n - 13 = |F_0|$ in $0LTQ_{n-1} - F_0$. Therefore, $N(V(P)) = F_0$ in $0LTQ_{n-1} - F_0$. Note that $2^{n-1} - (4n - 9) - 4 \geq 1$ for $n \geq 6$, by Lemma 3, then $LTQ_n[V(C_0) \cup V(C_1)]$ is connected. Then $LTQ_n - F$ just has two components, one of which is a 3-path.

If there exists a component $K_{1,3}$ in $0LTQ_{n-1} - F_0$, then $N_{0LTQ_{n-1}}(V(K_{1,3})) \subseteq F_0$. By Corollary 1, $|N(V(K_{1,3}))| \geq 4n - 13 = |F_0|$ in $0LTQ_{n-1} - F_0$. Therefore, $N(V(K_{1,3})) = F_0$ in $0LTQ_{n-1} - F_0$. Note that $2^{n-1} - (4n - 9) - 4 \geq 1$ for $n \geq 6$, by Lemma 3, $LTQ_n - F$ just has two components, one of which is a $K_{1,3}$.

If there exists a 4-cycle C in $0LTQ_{n-1} - F_0$, then

$N_{0LTQ_{n-1}}(C) \cap V(0LTQ_{n-1}) \subseteq F_0$. By Proposition 7,

$|N_{0LTQ_{n-1}}(V(C))| \geq 4(n - 1 - 2) = 4n - 12 > 4n - 13 = |F_0|$, a contradiction to $|F_0| = 4n - 13$. Therefore, $0LTQ_{n-1} - F_0$ has not a 4-cycle.

Case 5. $4n - 12 \leq |F_0| \leq 4n - 9$.

Since $4n - 12 \leq |F_0| \leq 4n - 9$ and $|F| \leq 4n - 9$, $|F_1| \leq (4n - 9) - (4n - 12) = 3$. By Lemma 1, $1LTQ_{n-1} - F_1$ is connected.

Suppose that $0LTQ_{n-1} - F_0$ is connected. Since $2^{n-1} - (4n - 9) \geq 1$, by Lemma 3, $LTQ_n - F$ is connected, a contradiction.

Suppose that $0LTQ_{n-1} - F_0$ is not connected. Let the components in $0LTQ_{n-1} - F_0$ be $G_1, G_2, \dots, G_k$ for $k \geq 2$ and $|V(G_1)| \leq |V(G_2)| \leq \dots \leq |V(G_k)|$. If $|V(G_r)| \geq 4(1 \leq r \leq k - 1)$, by Lemma 3, $|N(V(G_r)) \cap V(1LTQ_{n-1})| \geq 4$. If $k \geq 5$, by Lemma 3,

$|N(V(G_1)) \cup N(V(G_2)) \cup \dots \cup N(V(G_{k-1})) \cap V(1LTQ_{n-1})| \geq 4$. Combining this with $|F_1| \leq (4n - 9) - (4n - 12) = 3$, we have that $LTQ_n - F$ satisfies one of the following conditions:

- 1) $LTQ_n - F$ has four components, three of which are isolated vertices;
- 2) $LTQ_n - F$ has three components, one of which is isolated vertices and one of which is a K_2 ;
- 3) $LTQ_n - F$ has three components, two of which are isolated vertices;
- 4) $LTQ_n - F$ has two components, one of which is a path of length two;
- 5) $LTQ_n - F$ has two components, one of which is an isolated vertex;
- 6) $LTQ_n - F$ has two components, one of which is a K_2 ;
- 7) $LTQ_n - F$ is connected.

In this case, F is not a minimum 3-extra cut of LTQ_n , a contradiction.

4. The 3-Extra Diagnosability of the Locally Twisted Cube under the PMC Model

In this section, we shall show the 3-extra diagnosability of locally twisted cubes

under the PMC model.

Theorem 10. ([16] [22] [26]) *A system $G=(V, E)$ is g -extra t -diagnosable under the PMC model if and only if there is an edge $uv \in E$ with $u \in V \setminus (F_1 \cup F_2)$ and $v \in F_1 \Delta F_2$ for each distinct pair of g -extra faulty subsets F_1 and F_2 of V with $|F_1| \leq t$ and $|F_2| \leq t$.*

Lemma 11. *Let $n \geq 4$. Then the 3-extra diagnosability of the locally twisted cube LTQ_n under the PMC model is less than or equal to $4n-6$, i.e., $\tilde{t}_3(LTQ_n) \leq 4n-6$.*

Proof. Let A be defined in Lemma 7, and let $F_1 = N_{LTQ_n}(A)$, $F_2 = A \cup N_{LTQ_n}(A)$. By Lemma 7, $|F_1| = 4n-9$, $|F_2| = |A| + |F_1| = 4n-5$, $|V(LTQ_n[A])| \geq 4$ and $|V(LTQ_n - F_2)| \geq 4$, F_1 is a 3-extra cut of LTQ_n . Therefore, F_1 and F_2 are 3-extra faulty sets of LTQ_n with $|F_1| = 4n-9$ and $|F_2| = 4n-5$. Since $A = F_1 \Delta F_2$ and $N_{LTQ_n}(A) = F_1 \subset F_2$, there is no edge of LTQ_n between $V(LTQ_n) \setminus (F_1 \cup F_2)$ and $F_1 \Delta F_2$. By Theorem 10, we can deduce that LTQ_n is not 3-extra $(4n-5)$ -diagnosable under PMC model. Hence, by the definition of 3-extra diagnosability, we conclude that the 3-extra diagnosability of LTQ_n is less than $4n-5$, i.e., $\tilde{t}_3(LTQ_n) \leq 4n-6$.

Lemma 12. *Let $n \geq 5$. Then the 3-extra diagnosability of the locally twisted cube LTQ_n under the PMC model is more than or equal to $4n-6$, i.e., $\tilde{t}_3(LTQ_n) \geq 4n-6$.*

Proof. By the definition of 3-extra diagnosability, it is sufficient to show that LTQ_n is 3-extra $(4n-6)$ -diagnosable. By Theorem 10, to prove LTQ_n is 3-extra $(4n-6)$ -diagnosable, it is equivalent to prove that there is an edge $uv \in E(LTQ_n)$ with $u \in V(LTQ_n) \setminus (F_1 \cup F_2)$ and $v \in F_1 \Delta F_2$ for each distinct pair of 3-extra faulty subsets F_1 and F_2 of $V(LTQ_n)$ with $|F_1| \leq 4n-6$ and $|F_2| \leq 4n-6$.

Suppose, by way of contradiction, that there are two distinct 3-extra faulty subsets F_1 and F_2 of LTQ_n with $|F_1| \leq 4n-6$ and $|F_2| \leq 4n-6$, but the vertex set pair (F_1, F_2) is not satisfied with the condition in Theorem 10, i.e., there are no edges between $V(LTQ_n) \setminus (F_1 \cup F_2)$ and $F_1 \Delta F_2$. Without loss of generality, assume that $F_2 \setminus F_1 \neq \emptyset$.

Assume $V(LTQ_n) = F_1 \cup F_2$. Since $n \geq 5$, we have that $2^n = |V(LTQ_n)| = |F_1 \cup F_2| = |F_1| + |F_2| - |F_1 \cap F_2| \leq |F_1| + |F_2| \leq (4n-6) + (4n-6) = 8n-12$, a contradiction. Therefore, $V(LTQ_n) \neq F_1 \cup F_2$.

The following we discuss the case when $F_2 \setminus F_1 \neq \emptyset$ and $V(LTQ_n) \neq F_1 \cup F_2$.

Since there are no edges between $V(LTQ_n) \setminus (F_1 \cup F_2)$ and $F_1 \Delta F_2$, and F_1 is a 3-extra faulty set, $LTQ_n - F_1$ has two parts $LTQ_n - F_1 - F_2$ and $LTQ_n[F_2 \setminus F_1]$. Thus, every component G_i of $LTQ_n - F_1 - F_2$ satisfies $|V(G_i)| \geq 4$ and every component C_i of $LTQ_n[F_2 \setminus F_1]$ satisfies $|V(C_i)| \geq 4$. Similarly, every component C'_i of $LTQ_n[F_1 \setminus F_2]$ satisfies $|V(C'_i)| \geq 4$ when $F_1 \setminus F_2 \neq \emptyset$. Therefore, $F_1 \cap F_2$ is also a 3-extra faulty set. Since there are no edges between $V(LTQ_n - F_1 - F_2)$ and $F_1 \Delta F_2$, $F_1 \cap F_2$ is also a 3-extra cut. When $F_1 \setminus F_2 = \emptyset$, $F_1 \cap F_2 = F_1$ is also a 3-extra faulty set. Since there are no edges between $V(LTQ_n - F_1 - F_2)$ and $F_1 \Delta F_2$, $F_1 \cap F_2$ is a 3-extra cut. By

Theorem 8, $|F_1 \cap F_2| \geq 4n - 9$. Therefore,

$|F_2| = |F_2 \setminus F_1| + |F_1 \cap F_2| \geq 4 + 4n - 9 = 4n - 5$, which contradicts with that $|F_2| \leq 4n - 6$. So LTQ_n is 3-extra $(4n - 6)$ -diagnosable. By the definition of $\tilde{t}_3(LTQ_n)$, $\tilde{t}_3(LTQ_n) \geq 4n - 6$. The proof is complete.

Combining Lemmas 11 and 12, we have the following theorem.

Theorem 11. *Let $n \geq 5$. Then the 3-extra diagnosability of the locally twisted cubes LTQ_n under the PMC model is $4n - 6$.*

5. The 3-Extra Diagnosability of the Locally Twisted Cube under the MM* Model

Before discussing the 3-extra diagnosability of the locally twisted cube LTQ_n under the MM* model, we first give an existing result.

Theorem 12 ([3] [16] [26]) *A system $G = (V, E)$ is g -extra t -diagnosable under the MM* model if and only if for each distinct pair of g -extra faulty subsets F_1 and F_2 of V with $|F_1| \leq t$ and $|F_2| \leq t$ satisfies one of the following conditions.*

1) *There are two vertices $u, w \in V \setminus (F_1 \cup F_2)$ and there is a vertex $v \in F_1 \Delta F_2$ such that $uw \in E$ and $vw \in E$.*

2) *There are two vertices $u, v \in F_1 \setminus F_2$ and there is a vertex $w \in V \setminus (F_1 \cup F_2)$ such that $uw \in E$ and $vw \in E$.*

3) *There are two vertices $u, v \in F_2 \setminus F_1$ and there is a vertex $w \in V \setminus (F_1 \cup F_2)$ such that $uw \in E$ and $vw \in E$.*

Lemma 13. *Let $n \geq 4$. Then the 3-extra diagnosability of the locally twisted cube LTQ_n under the MM* model is less than or equal to $4n - 6$, i.e., $\tilde{t}_3(LTQ_n) \leq 4n - 6$.*

Proof. Let A be defined in Lemma 7, and let $F_1 = N_{LTQ_n}(A)$, $F_2 = A \cup N_{LTQ_n}(A)$. By Lemma 7, $|F_1| = 4n - 9$, $|F_2| = |A| + |F_1| = 4n - 5$, $|V(LTQ_n[A])| \geq 4$ and $|V(LTQ_n - F_2)| \geq 4$, F_1 is a 3-extra cut of LTQ_n . Therefore, F_1 and F_2 are 3-extra faulty sets of LTQ_n with $|F_1| = 4n - 9$ and $|F_2| = 4n - 5$. Since $A = F_1 \Delta F_2$ and $N_{LTQ_n}(A) = F_1 \subset F_2$, there is no edge of LTQ_n between $V(LTQ_n) \setminus (F_1 \cup F_2)$ and $F_1 \Delta F_2$. By Theorem 12, we can deduce that LTQ_n is not 3-extra $(4n - 5)$ -diagnosable under MM* model. Hence, by the definition of 3-extra diagnosability, we conclude that the 3-extra diagnosability of LTQ_n is less than $4n - 5$, i.e., $\tilde{t}_3(LTQ_n) \leq 4n - 6$.

A component of a graph G is odd or even according as it has an odd or even number of vertices. We denote by $o(G)$ the number of odd components of G .

Lemma 14. ([20]) *A graph $G = (V, E)$ has a perfect matching if and only if $o(G - S) \leq |S|$ for all $S \subseteq V$.*

Lemma 15. *Let $n \geq 7$. Then the 3-extra diagnosability of the locally twisted cube LTQ_n under the MM* model is more than or equal to $4n - 6$, i.e., $\tilde{t}_3(LTQ_n) \geq 4n - 6$.*

Proof. By the definition of the 3-extra diagnosability, it is sufficient to show that LTQ_n is 3-extra $(4n - 6)$ -diagnosable.

By Theorem 12, suppose, by way of contradiction, that there are two distinct

3-extra faulty subsets F_1 and F_2 of LTQ_n with $|F_1| \leq 4n-6$ and $|F_2| \leq 4n-6$, but the vertex set pair (F_1, F_2) is not satisfied with any one condition in Theorem 12. Without loss of generality, assume that $F_2 \setminus F_1 \neq \emptyset$. Similarly to the discussion on $V(LTQ_n) = F_1 \cup F_2$ in Lemma 12, we can deduce $V(LTQ_n) \neq F_1 \cup F_2$. Therefore, we have the following discussion for $V(LTQ_n) \neq F_1 \cup F_2$.

Claim 1. $LTQ_n - F_1 - F_2$ has no isolated vertex.

Suppose, by way of contradiction, that $LTQ_n - F_1 - F_2$ has at least one isolated vertex w . Since F_1 is a 3-extra faulty set, there is at least one vertex $u \in F_2 \setminus F_1$ such that u are adjacent to w . Since the vertex set pair (F_1, F_2) is not satisfied with any one condition in Theorem 12, by the condition (3) of Theorem 12, there is at most one vertex $u \in F_2 \setminus F_1$ such that u is adjacent to w . Therefore, there is just a vertex u is adjacent to w .

Case 1. $F_1 \setminus F_2 = \emptyset$.

If $F_1 \setminus F_2 = \emptyset$, then $F_1 \subseteq F_2$. Since F_2 is a 3-extra faulty set, every component G_i of $LTQ_n - F_1 - F_2$ has $|V(G_i)| \geq 4$. Thus, $LTQ_n - F_1 - F_2$ has no isolated vertex.

Case 2. $F_1 \setminus F_2 \neq \emptyset$.

Similarly, since $F_1 \setminus F_2 \neq \emptyset$, by the condition (2) of Theorem 12 and the hypothesis, we can deduce that there is just a vertex $v \in F_1 \setminus F_2$ such that v is adjacent to w .

Let $W \subseteq V(LTQ_n) \setminus (F_1 \cup F_2)$ be the set of isolated vertices in $LTQ_n[V(LTQ_n) \setminus (F_1 \cup F_2)]$, and H be the induced subgraph by the vertex set $V(LTQ_n) \setminus (F_1 \cup F_2 \cup W)$. Then for any $w \in W$, there are $(n-2)$ neighbors in $F_1 \cap F_2$. By Lemmas 14 and 3, $|W| \leq o(LTQ_n - (F_1 \cup F_2)) \leq |F_1 \cup F_2| = |F_1| + |F_2| - |F_1 \cap F_2| \leq (4n-6) + (4n-6) - (n-2) = 7n-10$. Assume $V(H) = \emptyset$. Then $2^n = |V(LTQ_n)| = |F_1 \cup F_2| + |W| = |F_1| + |F_2| - |F_1 \cap F_2| \leq (4n-6) + (4n-6) - (n-2) + (7n-10) = 14n-20$, a contradiction to that $n \geq 7$. So $V(H) \neq \emptyset$.

The following we discuss the case when $F_1 \setminus F_2 \neq \emptyset$, $F_2 \setminus F_1 \neq \emptyset$ and $V(H) \neq \emptyset$.

Since the vertex set pair (F_1, F_2) is not satisfied with the condition (1) of Theorem 12, and there are not isolated vertices in H , we induce that there is no edge between $V(H)$ and $F_1 \Delta F_2$. Note that $F_2 \setminus F_1 \neq \emptyset$. If $F_1 \cap F_2 = \emptyset$, then this is a contradiction to that LTQ_n is connected. Therefore, $F_1 \cap F_2 \neq \emptyset$. Thus, $F_1 \cap F_2$ is a vertex cut of LTQ_n . Since F_1 is a 3-extra faulty set of LTQ_n , we have that every component H_i of H has $|V(H_i)| \geq 4$ and every component C_i of $LTQ_n[W \cup (F_2 \setminus F_1)]$ has $|V(C_i)| \geq 4$. Since F_2 also is a 3-extra faulty set of LTQ_n , we have that every component C'_i of $LTQ_n[W \cup (F_1 \setminus F_2)]$ has $|V(C'_i)| \geq 4$. Note that $LTQ_n - (F_1 \cap F_2)$ has two parts: H and $LTQ_n[W \cup (F_1 \Delta F_2)]$. Let $b_i \in V(LTQ_n[W \cup (F_1 \Delta F_2)])$. If $b_i \in W$, then b_i has two neighbors $u \in V(C_i)$ and $v \in V(C'_i)$. Then $b_i \in V(C_i \cup C'_i)$ and $|V(C_i \cup C'_i)| \geq 4$. Thus, $F_1 \cap F_2$ is a 3-extra cut of LTQ_n . By Theorem 8, $|F_1 \cap F_2| \geq 4n-9$. Since $|V(C_i)| \geq 4$, $|F_2 \setminus F_1| \geq 3$. Since

$|F_1 \cap F_2| = |F_2| - |F_2 \setminus F_1| \leq (4n-6) - 3 = 4n-9$, we have $|F_1 \cap F_2| = 4n-9$. Then $|F_2 \setminus F_1| = 3$ and $|F_2| = 4n-6$. Similarly, $|F_1 \setminus F_2| = 3$, $|F_1| = 4n-6$. By Lemma 9, the locally twisted cube LTQ_n is tightly $(4n-9)$ super 3-extra connected, i.e., $LTQ_n - (F_1 \cap F_2)$ has two components, one of which is a subgraph of order 4. Noted that $|W| \leq 7n-10$. $2^n = |V(LTQ_n)| = |F_1 \setminus F_2| + |F_2 \setminus F_1| + |F_1 \cap F_2| + |V(H)| + |W| \leq 3+3+(4n-9)+4+(7n-10) = 11n-9$, a contradiction to $n \geq 7$. Therefore, $LTQ_n - F_1 - F_2$ has no isolated vertex when $F_1 \setminus F_2 \neq \emptyset$, $F_2 \setminus F_1 \neq \emptyset$ and $V(H) \neq \emptyset$. The proof of Claim 1 is complete.

Let $u \in V(LTQ_n) \setminus (F_1 \cup F_2)$. By Claim 1, u has at least one neighbor vertex in $LTQ_n - F_1 - F_2$. Since the vertex set pair (F_1, F_2) is not satisfied with any one condition in Theorem 12, by the condition (1) of Theorem 12, for any pair of adjacent vertices $u, w \in V(LTQ_n) \setminus (F_1 \cup F_2)$, there is no vertex $v \in F_1 \Delta F_2$ such that $uw \in E(LTQ_n)$ and $uv \in E(LTQ_n)$. It follows that u has no neighbor vertex in $F_1 \Delta F_2$. By the arbitrariness of u , there is no edge between $V(LTQ_n) \setminus (F_1 \cup F_2)$ and $F_1 \Delta F_2$. Since $F_2 \setminus F_1 \neq \emptyset$ and F_1 is a 3-extra faulty set, $|F_2 \setminus F_1| \geq 4$ and $|V(LTQ_n - F_2 - F_1)| \geq 4$. Since F_1 also is 3-extra faulty sets, $|F_1 \setminus F_2| \leq 4$ and $|V(LTQ_n - F_1 - F_2)| \geq 4$. Then $F_1 \cap F_2$ is a 3-extra cut of LTQ_n . By Theorem 8, we have $|F_1 \cap F_2| \geq 4n-9$. Therefore, $|F_2| = |F_2 \setminus F_1| + |F_1 \cap F_2| \geq 4 + (4n-9) = 4n-5$, which contradicts $|F_2| \leq 4n-6$. Therefore, LTQ_n is 3-extra $(4n-6)$ -diagnosable and $\tilde{t}_3(LTQ_n) \geq 4n-6$. The proof is complete.

Combining Lemmas 13 and 15, we have the following theorem.

Theorem 13. Let $n \geq 7$. Then the 3-extra diagnosability of the locally twisted cube LTQ_n under the MM model is $4n-6$.

Acknowledgements

This work is supported by the National Science Foundation of China (61370001).

References

- [1] Preparata, F.P., Metze, G. and Chien, R.T. (1967) On the Connection Assignment Problem of Diagnosable Systems. *IEEE Transactions on Computers*, **EC-16**, 848-854. <https://doi.org/10.1109/PGEC.1967.264748>
- [2] Maeng, J. and Malek, M. (1981) A Comparison Connection Assignment for Self-Diagnosis of Multiprocessor Systems. *Proceeding of 11th International Symposium on Fault-Tolerant Computing*, 173-175.
- [3] Sengupta, A. and Dahbura, A.T. (1992) On Self-Diagnosable Multiprocessor Systems: Diagnosis by the Comparison Approach. *IEEE Transactions on Computers*, **41**, 1386-1396. <https://doi.org/10.1109/12.177309>
- [4] Fàbrega, J. and Fiol, M.A. (1996) On the Extraconnectivity of Graphs. *Discrete Mathematics*, **155**, 49-57.
- [5] Gu, M.-M., Hao, R.-X. and Liu J.-B. (2017) On the Extraconnectivity of k -Ary n -Cube Networks. *International Journal of Computer Mathematics*, **94**, 95-106. <https://doi.org/10.1080/00207160.2015.1091070>
- [6] Chang, N.-W., Tsai, C.-Y. and Hsieh, S.-Y. (2014) On 3-Extra Connectivity and 3-Extra Edge Connectivity of Folded Hypercubes. *IEEE Transactions on Computers*, **63**, 1593-1599.

- [7] Zhang, M.-M. and Zhou, J.-X. (2015) On g -Extra Connectivity of Folded Hypercubes. *Theoretical Computer Science*, **593**, 146-153.
- [8] Hsieh, S.-Y. and Chang, Y.-H. (2012) Extraconnectivity of k -Ary n -Cube Networks. *Theoretical Computer Science*, **443**, 63-69.
- [9] Gu, M.-M. and Hao, R.-X. (2014) 3-Extra Connectivity of 3-Ary n -Cube Networks. *Information Processing Letters*, **114**, 486-491.
- [10] Lin, R.Z. and Zhang, H.P. (2016) The Restricted Edge-Connectivity and Restricted Connectivity of Augmented k -Ary n -Cubes. *International Journal of Computer Mathematics*, **93**, 1281-1298. <https://doi.org/10.1080/00207160.2015.1067690>
- [11] Lü, H. (2017) On Extra Connectivity and Extra Edge-Connectivity of Balanced Hypercubes. *International Journal of Computer Mathematics*, **94**, 813-820. <https://doi.org/10.1080/00207160.2016.1148813>
- [12] Hong, W.-S. and Hsieh, S.-Y. (2013) Extra Edge Connectivity of Hypercube-Like Networks. *International Journal of Parallel, Emergent and Distributed Systems*, **28**, 123-133. <https://doi.org/10.1080/17445760.2011.650696>
- [13] Xu, J.M., Wang, J.W. and Wang, W.W. (2010) On Super and Restricted Connectivity of Some Interconnection Networks. *Ars Combinatoria*, **94**, 1-8.
- [14] Peng, S.-L., Lin, C.-K., Tan, J.J.M. and Hsu, L.-H. (2012) The g -Good-Neighbor Conditional Diagnosability of Hypercube under PMC Model. *Applied Mathematics and Computation*, **218**, 10406-10412. <https://doi.org/10.1016/j.amc.2012.03.092>
- [15] Wang, S. and Han, W. (2016) The g -Good-Neighbor Conditional Diagnosability of n -Dimensional Hypercubes under the MM* Model. *Information Processing Letters*, **116**, 574-577. <https://doi.org/10.1016/j.ipl.2016.04.005>
- [16] Zhang, S. and Yang, W. (2016) The g -Extra Conditional Diagnosability and Sequential t/k -Diagnosability of Hypercubes. *International Journal of Computer Mathematics*, **93**, 482-497. <https://doi.org/10.1080/00207160.2015.1020796>
- [17] Ren, Y. and Wang, S. The Tightly Super 2-Extra Connectivity and 2-Extra Diagnosability of Locally Twisted Cubes. *Journal of Interconnection Networks* (To Appear).
- [18] Wang, S., Wang, Z. and Wang, M. (2016) The 2-Extra Connectivity and 2-Extra Diagnosability of Bubble-Sort Star Graph Networks. *The Computer Journal*, **59**, 1839-1856. <https://doi.org/10.1093/comjnl/bxw037>
- [19] Wang, S. and Yang, Y. (2017) The 2-Good-Neighbor (2-Extra) Diagnosability of Alternating Group Graph Networks under the PMC Model and MM* Model. *Applied Mathematics and Computation*, **305**, 241-250. <https://doi.org/10.1016/j.amc.2017.02.006>
- [20] Bondy, J.A. and Murty, U.S.R. (2007) Graph Theory. Springer, New York.
- [21] Ren, Y. and Wang, S. (2016) Some Properties of the g -Good-Neighbor (g -Extra) Diagnosability of a Multiprocessor System. *American Journal of Computational Mathematics*, **6**, 259-266. <https://doi.org/10.4236/ajcm.2016.63027>
- [22] Dahbura, A.T. and Masson, G.M. (1984) An $O(n^{2.5})$ Fault Identification Algorithm for Diagnosable Systems. *IEEE Transactions on Computers*, **33**, 486-492. <https://doi.org/10.1109/TC.1984.1676472>
- [23] Fan, J. (2002) Diagnosability of Crossed Cubes under the Comparison Diagnosis Model. *IEEE Transactions on Parallel and Distributed Systems*, **13**, 1099-1104. <https://doi.org/10.1109/TPDS.2002.1041887>
- [24] Fan, J., Zhang, S., Jia, X. and Zhang, G. (2009) The Restricted Connectivity of Locally Twisted Cubes. *10th International Symposium on Pervasive Systems, Algorithms, and Networks (ISPAN)*, Kaohsiung, 14-16 December 2009, 574-578.

- <https://doi.org/10.1109/I-SPAN.2009.48>
- [25] Lai, P.-L., Tan, J.J.M., Chang, C.-P. and Hsu, L.-H. (2005) Conditional Diagnosability Measures for Large Multiprocessor Systems. *IEEE Transactions on Computers*, **54**, 165-175. <https://doi.org/10.1109/TC.2005.19>
 - [26] Yuan, J., Liu, A., Ma, X., Liu, X., Qin, X. and Zhang, J. (2015) The g -Good-Neighbor Conditional Diagnosability of k -Ary n -Cubes under the PMC model and MM* Model. *IEEE Transactions on Parallel and Distributed Systems*, **26**, 1165-1177. <https://doi.org/10.1109/TPDS.2014.2318305>
 - [27] Yang, X., Evans, D.J. and Megson, G.M. (2005) The Locally Twisted Cubes. *International Journal of Computer Mathematics*, **82**, 401-413. <https://doi.org/10.1080/0020716042000301752>
 - [28] Ren, Y. and Wang, S. The 1-Good-Neighbor Connectivity and Diagnosability of Locally Twisted Cubes. *Chinese Quarterly Journal of Mathematics* (To Appear).
 - [29] Feng, R., Bian, G. and Wang, X. (2011) Conditional Diagnosability of the Locally Twisted Cubes under the PMC Model. *Communications and Network*, **3**, 220-224. <https://doi.org/10.4236/cn.2011.34025>
 - [30] Hsieh, S.-Y., Huang, H.-W. and Lee, C.-W. (2016) {2,3}-Restricted Connectivity of Locally Twisted Cubes. *Theoretical Computer Science*, **615**, 78-90. <https://doi.org/10.1016/j.tcs.2015.11.050>
 - [31] Zhu, Q., Wang, X.-K. and Cheng, G. (2013) Reliability Evaluation of BC Networks. *IEEE Transactions on Computers*, **62**, 2337-2340. <https://doi.org/10.1109/TC.2012.106>



Scientific Research Publishing

Submit or recommend next manuscript to SCIRP and we will provide best service for you:

Accepting pre-submission inquiries through Email, Facebook, LinkedIn, Twitter, etc.
 A wide selection of journals (inclusive of 9 subjects, more than 200 journals)
 Providing 24-hour high-quality service
 User-friendly online submission system
 Fair and swift peer-review system
 Efficient typesetting and proofreading procedure
 Display of the result of downloads and visits, as well as the number of cited articles
 Maximum dissemination of your research work

Submit your manuscript at: <http://papersubmission.scirp.org/>

Or contact ajcm@scirp.org

